

Introduction to Concurrency

SWEN-342

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What Is a Process?

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What Is a Thread?

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A process has a self-contained execution environment. A process generally has a complete, private set of basic run-time resources; in particular, each process has its own memory space.

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What Is a Thread?

Threads exist within a process — every process has at least one. Threads share the process's resources, including memory and open files. This makes for efficient, but potentially problematic, communication.

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 - Better resource utilization.
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- Original (OS centric processes)
 - Better resource utilization.
 - Fairness among multiple users with multiple computations.
- Current (process centric threads)
 - Exploiting multiple processors
 - [Moore's Law](#) running out of steam (multi-core).
 - Modeling: Divide & conquer on loosely related tasks.
 - Simplify handling asynchronicity (e.g., mouse events)
 - Throughput (even on single CPU systems)
 - Responsiveness

The Perils of Concurrency

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 - Race conditions
 - Memory barrier (caching)
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- Performance
 - TANSTAAFL (There Ain't No Such Thing As A Free Lunch)
 - Context switching overhead
 - Disabled compiler optimizations

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- Testing, hair-pulling, and Heisenbugs

The Perils of Concurrency

- **1985-1987 -- Therac-25 medical accelerator.** A radiation therapy device malfunctions and delivers lethal radiation doses at several medical facilities. Based upon a previous design, the [Therac-25](#) was an "improved" therapy system that could deliver two different kinds of radiation: either a low-power electron beam (beta particles) or X-rays. The Therac-25's X-rays were generated by smashing high-power electrons into a metal target positioned between the electron gun and the patient. A second "improvement" was the replacement of the older Therac-20's electromechanical safety interlocks with software control, a decision made because software was perceived to be more reliable.
- What engineers didn't know was that both the 20 and the 25 were built upon an operating system that had been kludged together by a programmer with no formal training. Because of a subtle bug called a "[race condition](#)," a quick-fingered typist could accidentally configure the Therac-25 so the electron beam would fire in high-power mode but with the metal X-ray target out of position. At least five patients die; others are seriously injured.

(Source: "History's Worst Software Bugs", Wired)

The Ultimate Culprit - Shared, Mutable State

- Most of your development has been in imperative languages.
- The fundamental operation is assignment to change state.
 - Assignable variables are mutable.
 - May be exposed as public (bad karma).
 - May be exposed via interface methods (medium warm karma).
 - Things get tricky very fast when > 1 thread can invoke a mutating function.

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```
public class Counter {  
    private int count = 0 ;  
  
    public void increment() {  
        count = count + 1 ;  
    }  
  
    public int getCount() {  
        return count ;  
    }  
}
```

If we call **increment()** 10,000 times
and then call **getCount()**, what value is
returned?

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- Three basic approaches:
 - Make things immutable.
 - Hide shared state behind sequential access.
 - Provide mechanisms to support controlled access to shared, mutable state.

Other Issues

- Thread management
 - How many threads at one time?
 - Allocation of tasks to threads.
 - Thread scheduling.
- Higher level constructs
 - Fork / join
 - Callables & Futures
- Distributed state management
 - State consistency
 - Decision consensus